

Course outline

How does an NPTEL online course work?

Week 0

Week 1

Week 2

• Lecture 05: Hierarchy Theorems

• Lecture 06: Introduction to Space Complexity

• Lecture 07: Savitch's Theorem

 Quiz : Assignment 2

• Feedback for Week 2

• Assignment 2 Solution

Week 3

Week 4

week 5

week 6

Week 7

Week 8

week 9

Week 10

Week 11

Week 12

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Assignment 2

The due date for submitting this assignment has passed.

Due on 2021-02-07, 23:59 IST.

As per our records you have not submitted this assignment.

1) Mark all the correct options

6 points

$$NP = SPACE(n) \implies SPACE(n) = SPACE(n^2)$$

$$SPACE(n) \neq SPACE(n^2) \implies P \neq SPACE(n)$$

 For time constructible functions t_1, t_2, g ,

$$TIME(t_1(n)) \subsetneq TIME(t_2(n)) \implies TIME(t_1(g(n))) \subsetneq TIME(t_2(g(n)))$$

 For space constructible functions s_1, s_2, f ,

$$NSPACE(s_1(n)) \subseteq SPACE(s_2(n)) \implies NSPACE(s_1(f(n))) \subseteq SPACE(s_2(f(n)))$$

No, the answer is incorrect.
Score: 0
Accepted Answers:

$$NP = SPACE(n) \implies SPACE(n) = SPACE(n^2)$$

$$SPACE(n) \neq SPACE(n^2) \implies P \neq SPACE(n)$$

 For space constructible functions s_1, s_2, f ,

$$NSPACE(s_1(n)) \subseteq SPACE(s_2(n)) \implies NSPACE(s_1(f(n))) \subseteq SPACE(s_2(f(n)))$$

2) Which of the following statements is/are true?

6 points

 If f is time constructible, then

$$TIME(f(n)) \subsetneq NTIME(f(n))$$

 There exists a function f such that $TIME(f(n)) = TIME(2^{f(n)})$

 If f is a time constructible function, then

$$TIME\left(o\left(\frac{f(n)}{\log f(n)}\right)\right) \subsetneq TIME(f(n))$$

 The class P could in principle collapse to $TIME(n^k)$ for some fixed k
No, the answer is incorrect.
Score: 0
Accepted Answers:

 There exists a function f such that $TIME(f(n)) = TIME(2^{f(n)})$

 If f is a time constructible function, then

$$TIME\left(o\left(\frac{f(n)}{\log f(n)}\right)\right) \subsetneq TIME(f(n))$$

3) Mark all the true statements

4 points

$$P \subsetneq P^P$$

 For any oracle A , $P^A \subsetneq EXP^A$

 Let $TIME^A(f(n)) = \{L \mid \text{there is an oracle machine } M^A \text{ that decides } L \text{ and runs in time } f(n)\}$

 There exists an oracle A such that $TIME^A(n^2) = TIME^A(n^3)$

 There exists an oracle A relative to which P and NP are different

No, the answer is incorrect.
Score: 0
Accepted Answers:

 For any oracle A , $P^A \subsetneq EXP^A$

 There exists an oracle A relative to which P and NP are different

4) Which of the following statements is/are true?

2 points

$$TIME(f(n)) \subseteq SPACE(f(n))$$

$$NSPACE(n^2 + 3) \subseteq SPACE(n^4 + 6n^2 + 9)$$

$$NSPACE(f(n)) \subseteq SPACE(f(n))$$

$$NPSPACE \subseteq EXP$$

No, the answer is incorrect.
Score: 0
Accepted Answers:

$$TIME(f(n)) \subseteq SPACE(f(n))$$

$$NSPACE(n^2 + 3) \subseteq SPACE(n^4 + 6n^2 + 9)$$

$$NPSPACE \subseteq EXP$$

5) Mark all the statements that are known to be true

2 points

$$NL = P$$

$$EXPSPACE = NEXPSPACE$$

 Directed graph reachability problem can be solved in $O(\log(n))$ space

 Directed graph reachability problem can be solved in $O(\log^2(n))$ space

No, the answer is incorrect.
Score: 0
Accepted Answers:

$$EXPSPACE = NEXPSPACE$$

 Directed graph reachability problem can be solved in $O(\log^2(n))$ space

 6) Mark all the true statements about the polynomial time reduction relation \leq_p

2 points

 \leq_p is reflexive

 \leq_p is symmetric

 \leq_p is transitive

 \leq_p is one-to-one

No, the answer is incorrect.
Score: 0
Accepted Answers:
 \leq_p is reflexive

 \leq_p is transitive

 7) Given a DFA D and an input x , we need to decide if D accepts x .

2 points

What is the smallest class in which we can do this?

 L

 NL

 P

 NP
No, the answer is incorrect.
Score: 0
Accepted Answers:
 L

 8) The class NL is closed under which of the following operations? Mark all that apply

2 points

 Intersection

 Union

 Concatenation

 Kleene Closure

No, the answer is incorrect.
Score: 0
Accepted Answers:

Intersection

Union

Concatenation

Kleene Closure

 9) For any positive integer n , let n' be the integer whose binary representation is the reverse of the binary representation of n (assume that binary representation of n has no leading zeros). Let $f(n) = n + n'$

4 points

 for all n . Consider the following statements:

$$L_1 = \{\langle m, n \rangle \mid f(m) = n\}$$

$$L_2 = \{\langle m, n \rangle \mid f(f(m)) = n\}$$

 Only L_1 is in L

 Only L_2 is in L

 Both L_1 and L_2 are in L

 Both L_1 and L_2 are not in L
No, the answer is incorrect.
Score: 0
Accepted Answers:

 Both L_1 and L_2 are in L